

The Curry-Howard isomorphism for dummies

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PPS/ πr^2

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All you ever wanted to know and were afraid to ask!
(or did not care about)

In this talk

- A gentle introduction to the Curry-Howard isomorphism
- Bird's eye view, no technical stuff
- A bit of advertising

A brief, biased history of Logic: The Dark Ages

- 4th c. **B.C.** The greek philosopher Aristotle writes the *Organon*.

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Meanwhile, heated discussion between byzantine scholars about the sex of angels.

- **1901.** Russel shows that there is no set of all sets.

"The 20th century is the century of mathematics."

A brief, biased history of Logic: The Age of Enlightenment

From the 30's onwards, foundations of mathematics and computer science become intertwined.

Formal logic

- **1929.** Completeness theorem
- **1931.** Incompleteness theorem
- **1936.** Undefinability theorem

Computer science

- **193X.** λ -calculus
- **1936.** Turing Machines
- **1936.** Halting problem

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... and there is a good reason for that.

Through the looking glass

The mathematician

Theorem. For all $n \in \mathbb{N}$, there exists $p \in \mathbb{N}$ such that $n = 2p$ or $n = 2p + 1$.

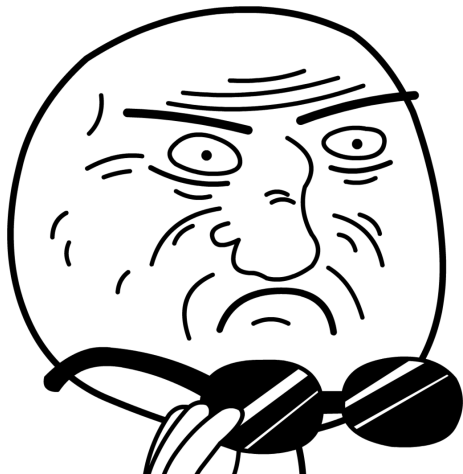
Proof. By induction on n .

- If $n = 0$ then this is obvious.
- Otherwise, assume that $n = m + 1$. By the induction hypothesis, we know that there exists some p such that $m = 2p$ or $m = 2p + 1$.
 - In the first case, $n = 2p + 1$.
 - Otherwise $n = 2(p + 1)$.

The programmer

```
val div2 : int -> int * bool
(* [div2 n] returns the integer
division by 2 of [n] together with
a boolean indicating if [n] is
even. *)
```

```
let rec div2 n = match n with
| 0 -> (0, true)
| m + 1 ->
  let (p, even) = div2 m in
  if even then (p, false)
  else (p + 1, true)
```

That's the same thing.

The Curry-Howard isomorphism

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- a.k.a. the “proof-as-program” correspondence
- Proofs compute! ♡

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Furthermore, the statement of a theorem correspond to the type of the corresponding program.

- We need our programming language to be rich enough
- Programming languages in the wild have crappy type systems (if any)

“Do not try to do maths with C++/Java/Python at home, kids.”

A Rosetta's Stone

Logic

Proofs

Formula

CS

Programs

Types

A Rosetta's Stone

Logic

Proofs

Formula

A implies B

A and B

A or B

falsity

truth

for all $x \in A$, $B(x)$

CS

Programs

Types

function from A to B

pair of A and B

tagged union of A and B

empty type

singleton type

dependent product from A to B

A Rosetta's Stone

Logic

Proofs

Formula

A implies B

A and B

A or B

falsity

truth

for all $x \in A$, $B(x)$

Axiom

Soundness theorem

Completeness theorem

Incompleteness theorem

CS

Programs

Types

function from A to B

pair of A and B

tagged union of A and B

empty type

singleton type

dependent product from A to B

System primitive

Compiler

Debugger

Infinite loop

A paradigmatic shift

Proofs are relevant objects.

“Any list can be quotiented by permutations.”

- Quicksort
- Mergesort
- Bogosort

A paradigmatic shift

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Proofs are first-class objects.

$\forall(A : \mathbf{Set}). \forall(x : A). \forall(p\ q : x = x). p = q$ (?)

Which logic for which programs?

Standard members of each community will complain.

The mathematician

"This logic is crappy, it does not feature the following principles I am acquainted with."

Either A or not A hold.

Every bounded monotone sequence has a limit.

Two sets with same elements are equal.

Every formula is equivalent to its prenex form.

...

The programmer

"This language is crappy, it does not feature the following structures I am acquainted with."

```
printf("Hello world")
  x <- 42
  while true { ... }
goto #considered_harmful
  fork()
  ...
```

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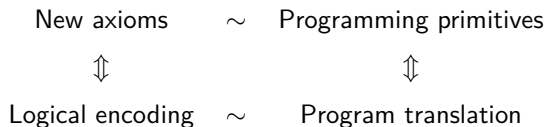
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INTUITIONISTIC LOGIC

FUNCTIONAL LANGUAGE

What Curry-Howard taketh, it giveth back.

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The prototypical example: Classical logic

We can implement classical logic through this scheme.

New axiom

$$A \vee \neg A$$

Used by your next-door mathematician



Logical encoding

double-negation translation

Gödel already did that by the 30's

~

Programming primitives

`callcc`

From Scheme, though a bit arcane



Program translation

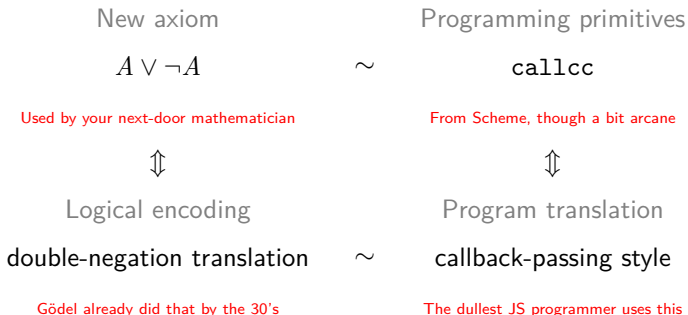
callback-passing style

The dullest JS programmer uses this

~

The prototypical example: Classical logic

We can implement classical logic through this scheme.



What does it do, computationally?

A little tale



- Did the Devil lie?

Morale

- Did the Devil lie?
- That's not clear!
- Ensure you agree on the semantics before bargaining with the Devil.
- Hint: that's a parable.

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Program translations can change the expected behaviour

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Program translations can change the expected behaviour

...and it can invalidate some other principles

- Axiom of choice is trivial in intuitionistic logic
- Axiom of choice is a monster from outer space in classical logic

An ever-growing field

The back-and-forth journey between world is fruitful.

- Double negation \sim first-class callbacks
- Friedman's translation \sim dynamic exceptions
- Cohen's forcing \sim global variables
- Dialectica translation \sim scoped `gotos` (*à la* Python's `yield`)

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...and there is more

- *Trending topic*: HoTT
- Proofs as paths, types as homotopy spaces
- Discovered by Voevodsky (2002 Fields medal)

WARNING

In the following next slides, you are going to be exposed to blatant advertising material. You have been advised.

The *fine fleur* of the Curry-Howard isomorphism.

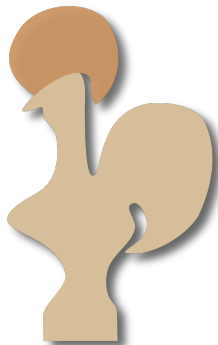
- **Your 2-in-1 proof assistant!**

- A programming language
- A mathematical system

- Developed by πr^2 team
- Est. 1984 (Coquand, Huet, Paulin...)
- 2013 ACM Software System Award

Previously: LLVM, Unix, TCP/IP, TeX, Apache, Java...

- Version 8.5 beta just released.



A success story

Coq is increasingly used in **The Real World**®.

Maths

- Four-Colour Theorem
- Feit-Thomson Theorem
- HoTT/Coq

CS

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- Bedrock
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Try it out today!

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 - Logic and computation unified
 - A theory of Everything

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- Food for thought
 - Why aren't mathematics closed-source?
 - What is a mathematical bug?
 - Is a iTheorem really cooler than a WinTheorem32?
 - Do Python coders truly believe that $2 + 2 = 5$?

Any question?