

Hurewicz fibrations in elementary toposes

Revised version

BY KRZYSZTOF WORYTKIEWICZ

Université de Savoie Mont-Blanc

Realizability Boreal Summer School, Piriapolis, July 19, 2016

Intro

Question. *What is a model of HoTT? What is a model of HoTT with univalence?*

Remark. We won't address univalence in this talk...

Typoi

Definition. Let \mathbb{C} be a category with finite limits.

1. A class of maps $\mathcal{F} \subseteq \mathbb{C}_1$ is a class of fibrations provided it contains all isos and is stable under composition as well as under base change.
2. Let $\mathcal{F} \subseteq \mathbb{C}_1$ be a class of fibrations. An object $X \in \mathbb{C}$ is \mathcal{F} -fibrant if $!_X: X \rightarrow 1$ is a fibration.
3. Let $\mathcal{F} \subseteq \mathbb{C}_1$ be a class of fibrations. $(\mathbb{C}, \mathcal{F})$ is a tribe provided every object $X \in \mathbb{C}$ is \mathcal{F} -fibrant.

Definition. Let \mathbb{C} be a category with pullbacks, $f: A \rightarrow B$ be a map and $(E, p) \in \mathbb{C}/A$. An object $\prod_f (E, p) = (\prod_f (E), \prod_f (p)) \in \mathbb{C}/B$ equipped with a morphism $\epsilon: \prod_f (E) \times_B A \rightarrow E$ is the product of E along f if the morphism

$$\epsilon! f^*(-): \mathbb{C}/B \left((X, u), \prod_f (E, p) \right) \rightarrow \mathbb{C}/B(f^*(X, u), (E, p))$$

is iso for every $(X, u) \in \mathbb{C}/B$. The map ϵ is called evaluation.

Definition. A tribe $(\mathbb{C}, \mathcal{F})$ is \prod -closed provided

- i. every fibration $p: E \rightarrow A$ has a product along every fibration $f: A \rightarrow B$;
- ii. and $\prod_f (E, p)$ is again a fibration.

Example.

1. Any LCCC with all the maps.
2. Any CCC with projections.
3. Small groupoids with Grothendieck fibrations (**Hoffmann-Streicher**) .
4. Kan complexes with Kan fibrations (**Streicher, Voevodsky**) .
5. Type theory terms with display maps (**Gambino-Garner**)

Definition. Let $(\mathbb{C}, \mathcal{F})$ be a tribe.

1. A morphism in \mathbb{C} is anodyne provided $c \in {}^{\mathfrak{h}}\mathcal{F}$.
2. Let $\mathcal{A} \subseteq \mathbb{C}_1$ be the class of anodyne morphisms. The tribe $(\mathbb{C}, \mathcal{F})$ is homotopical provided
 - i. $(\mathcal{A}, \mathcal{F})$ is a factorisation system;
 - ii. anodyne morphisms are stable under base change along a fibration.

Remark. (Joyal) We have $(i) \Rightarrow (ii)$ provided the tribe is \square - closed.

Definition. Let $(\mathbb{C}, \mathcal{F})$ be a homotopical tribe and $A \in \mathbb{C}$. A path object $\mathbf{P}A$ is given by an $(\mathcal{A}, \mathcal{F})$ -factorisation of the diagonal $\Delta: A \rightarrow A \times A$. A homotopy with respect to a path object is called path homotopy.

Remark. Let $(\mathbb{C}, \mathcal{F})$ be a homotopical tribe. A path objects exists and can be “lifted” to slices.

Theorem. (Joyal) Let $(\mathbb{C}, \mathcal{F})$ be a homotopical tribe. The path homotopy relation is a congruence on \mathbb{C} .

Definition. A tribe $(\mathbb{C}, \mathcal{F})$ is a typos provided

i. it is homotopical and \prod - closed;

ii. the product functor \prod_f preserves the path homotopy relation for every fibration f .

Theorem. (Hoffman-Streicher) *The tribe of small groupoids and Grothendieck fibrations is a typos.*

Theorem. (Awodey-Warren-Voevodsky) *The tribe of Kan complexes and Kan fibrations is a typos.*

Theorem. (Gambino-Garner) *The tribe of type theory terms and display maps is a typos.*

Question. *How about realisability toposes?*

Remark. What follows is a vast generalisation of the material in Jaap van Oosten's seminal paper *Notion of Homotopy for the Effective Topos* (2010).

Intervals

Fix an elementary topos \mathbb{T} .

Definition.

1. $X \in \mathbb{T}$ is connected if $\mathbb{T}(X, 1 + 1) = \{\text{inl} \circ !_X, \text{inr} \circ !_X\}$.
2. $I \in \mathbb{T}$ is an elementary interval if it is connected and has precisely two distinct global elements $\partial_0, \partial_1: 1 \rightarrow I$.
3. $I_n \stackrel{\text{def.}}{=} \underbrace{I +_{\partial_0, \partial_1} \cdots +_{\partial_0, \partial_1} I}_{n \times}$.

Remark.

1. $I_0 = 1$ and $I_1 = I$.
2. I_n has precisely $n + 1$ global elements $\#i_n: 1 \rightarrow I_n$ corresponding to the injections into the defining wide pushout.
3. There is the obvious linear order on $\mathbb{T}(1, I_n) = \{\#0, \dots, \#n\}$.
4. I_n is connected for all $n \in \mathbb{N}$.

Definition. A Hurewicz topos is an elementary topos with NNO, equipped with a distinguished elementary interval.

Remark. In a Hurewicz topos \mathbb{T} coproducts $\coprod_{n \in \mathbb{N}} X^{I_n}$ exist for all $X \in \mathbb{T}$ since general bounded (co)limits exist in any topos.

Example. Any Grothendieck topos.

Example. The *effective topos* **Eff** where the distinguished elementary interval is the assembly

$$I = (\{0, 1\}; E(0) = \{0, 1\}, E(1) = \{1, 2\})$$

since there is no uniform realizer for the map

$$\begin{array}{ccc} e: I & \longrightarrow & 1 + 1 \\ i & \mapsto & i \end{array}$$

where $i \in \{0, 1\}$. Notice that in (\mathbf{Eff}, I) an object (X, \approx) is connected provided $E(x) \cap E(x') \neq \emptyset$ for all $x, x' \in X$.

Example. Any realizability topos over a PCA.

Question. *Any realisability topos?*

Remark. As pointed out by several people in the audience, classical toposes (e.g. realizability toposes over *Krivine structures*) cannot be Hurewicz since internal Excluded Middle stands in the way...

Paths in Hurewicz toposes

Assume \mathbb{T} is Hurewicz.

Definition.

1. $s_n \stackrel{\text{def.}}{=} \#0: 1 \rightarrow I_n$ and $t_n \stackrel{\text{def.}}{=} \#n \rightarrow I_n$ are called I_n 's endpoints.
2. A map $f: I_m \rightarrow I_n$ is
 - endpoint-preserving if $f \circ s_m = s_n$ and $f \circ t_m = t_n$;
 - order-preserving if $\#i_m \leq \#j_m \Rightarrow f \circ \#i_m \leq f \circ \#j_m$.
3. A order and endpoint preserving map is called degeneracy.

Remark. Suppose $k: I_m \rightarrow I_n$ is a degeneracy. Then k is epi and $m \geq n$.

Definition. Let $X \in \mathbb{T}$. The path object $X^{\langle I \rangle}$ of X is given by

$$X^{\langle I \rangle} \stackrel{\text{def.}}{=} \coprod_{n \in \mathbb{N}} X^{I_n} / \sim$$

where $\sigma \sim \theta$ if there is a degeneracy δ such that $\sigma = \theta \circ \delta$ or $\theta = \sigma \circ \delta$.

Remark.

1. Any path $[\sigma] \in X^{\langle I \rangle}$ has a canonical representative of minimal length.
2. For any two paths $[\sigma], [\theta] \in X^{\langle I \rangle}$ there are always representatives of same length.
3. Let $\rho: I_m \rightarrow X$ and $\rho': I_n \rightarrow X$ be representatives of a path $[\sigma] \in X^{\langle I \rangle}$. Then

$$\rho \circ \#m = \rho' \circ \#n$$

(since ρ and ρ' are related by a degeneracy).

4. In a Hurewicz topos, *path-connectedness* and *connectedness* are equivalent notions.

Definition. A constant path is a path with the path of length 0 among its representatives.

Notation. We shall write $[\sigma_m]$ if there is a need to insist that the representative's domain is I_m .

Remark. The source and target maps $s, t: X^{\langle I \rangle} \rightarrow X$ given by evaluations $s([\sigma_m]) = \sigma_m(0)$ and $t([\sigma_m]) = \sigma_m(m)$ respectively determine an internal graph $X^{\langle I \rangle} \rightrightarrows X$ in \mathbb{T} , since representatives differ by a degeneracy.

Notation. We shall abuse notation and write $X^{\langle I \rangle}$ for the internal graph $X^{\langle I \rangle} \rightrightarrows X$.

Proposition. The internal graph $X^{\langle I \rangle}$ is an internal category with

i. identity $c: X \rightarrow X^{\langle I \rangle}$ given by $c(x) = \left[1 \xrightarrow{[x]} X \right]$;

ii. composition $(- * -): X^{\langle I \rangle} \times_X X^{\langle I \rangle} \rightarrow X^{\langle I \rangle}$ given by

$$([\sigma_m] * [\theta_n])(i) = [\mathbf{if } i \leq m \mathbf{ then } \sigma_m(i) \mathbf{ else } \theta_n(i)]$$

Moreover, there is a contravariant involution $(-)^{\text{rev}}: X^{\langle I \rangle} \rightarrow X^{\langle I \rangle}$ which is constant on objects and given by

$$[\sigma_n]^{\text{rev}}(i) = \sigma_n(n - i)$$

on maps.

Remark. $(-)^{\langle I \rangle}: \mathbb{T} \rightarrow \mathbb{T}$ is a functor acting on paths by postcomposition:

$$f^{\langle I \rangle}([\sigma]) = [f \circ \sigma]$$

Moreover, all the associated maps are natural.

Remark. There is a factorisation of the diagonal map

$$\begin{array}{ccc}
 & X \langle I \rangle & \\
 c \nearrow & & \searrow (s,t) \\
 X & \xrightarrow{\Delta_X} & X \times X
 \end{array}$$

Definition. Let $f, g: X \rightarrow Y$ be maps. A homotopy $H: f \rightsquigarrow g$ from f to g is given by the commuting diagram

$$\begin{array}{ccccc}
 & & X & & \\
 & & \downarrow & & \\
 f \swarrow & & H & & \searrow g \\
 Y & \xleftarrow{s} & X \langle I \rangle & \xrightarrow{t} & Y
 \end{array}$$

H is constant on a subobject $X' \triangleleft X$ if $H(x) = c(x)$ for all $x \in X'$.

Remark. So for any $x \in X$ we have a path $H(x): f(x) \rightsquigarrow g(x)$.

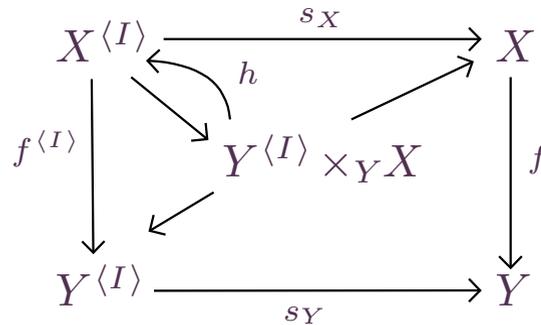
Homotopy equivalences

Definition. A homotopy equivalence is a map $u: X \rightarrow Y$ which has an up-to-homotopy inverse $v: Y \rightarrow X$.

Remark. The map v is a homotopy equivalence as well, called *the inverse homotopy equivalence*.

Fibrations

Definition. A section h of the map $(f^{(I)}, s_X)$



is called Hurewicz connection. A map which admits a Hurewicz connection is called Hurewicz fibration.

Notation. We shall write \mathcal{H} for the class of Hurewicz fibrations.

Remark. A Hurewicz fibration $f: X \rightarrow Y$ is thus a map with a path lifting property: for any path $\sigma: y \rightsquigarrow y'$ in Y and any $x \in X$ such that $f(x) = y$ there is a path θ in X such that $f \circ \theta = \sigma$:

$$\begin{array}{ccc}
 X & \xrightarrow[\exists]{x \text{ --- } \rightarrow x'} & \\
 f \downarrow & \text{-----} & \\
 Y & \xrightarrow{y \text{ --- } \rightarrow y'} &
 \end{array}$$

Definition. $X \in \mathbb{T}$ is fibrant if $!_X: X \rightarrow 1$ is a fibration.

Proposition.

1. Fibrations are stable under pullback and composition.
2. Any iso is a fibration.
3. Any object is fibrant.
4. $(s, t): X^{\langle I \rangle} \rightarrow X$ is a fibration for any $X \in \mathbb{T}$.

Strong deformation retracts

Definition. $X \in \mathbb{T}$ is a strong deformation retract of $Y \in \mathbb{T}$ if there is a map $e: X \rightarrow Y$ admitting a retraction $r: Y \rightarrow X$ such that there is a homotopy $H: \text{id}_Y \rightsquigarrow e \circ r$ constant on X (that is $H(x): x \rightsquigarrow (e \circ r)(x)$ is the constant path $c(x)$ for all $x \in X$). The split mono e is called sdr-insertion.

Remark.

1. Any sdr-insertion is a homotopy equivalence.
2. Sdr-insertions are stable under pullback along a fibration.

Factorisations

Definition. A map $a \in {}^{\mathfrak{h}}\mathcal{H}$ is called anodyne.

Proposition. Sdr-insertions are anodyne.

Definition. Let $f: X \rightarrow Y$ be a map. The object M_f given by the pullback

$$\begin{array}{ccc} M_f & \xrightarrow{p_2} & Y \langle I \rangle \\ \downarrow p_1 & \lrcorner & \downarrow s \\ X & \xrightarrow{f} & Y \end{array}$$

is called f 's mapping track.

Remark. Pulling back the factorisation $\Delta_Y = (s, t) \circ c$ of the diagonal yields a factorisation of the diagonal $\Delta_X = p_1 \circ c'$ with $p_1 \in \mathcal{H}$ and $c' \in {}^{\mathfrak{m}}\mathcal{H}$:

$$\begin{array}{ccc}
 & X & \xrightarrow{p'_2} & Y \\
 & \downarrow c' \lrcorner & & \downarrow c \\
 \Delta_X & M_{f \times f} & \xrightarrow{p_2} & Y \langle I \rangle \\
 & \downarrow p_1 \lrcorner & & \downarrow (s, t) \\
 & X \times X & \xrightarrow{f \times f} & Y \times Y \\
 & & & \Delta_Y
 \end{array}$$

Theorem. Any map $f: X \rightarrow Y$ factors through the mapping track M_f as $f = h \circ a$ with a anodyne and h a Hurewicz fibration.

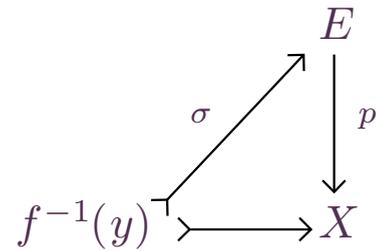
Products along maps

Remark. It is well-known that in any topos \mathbb{T} and any map $f: X \rightarrow Y$ in \mathbb{T} , the pullback functor $f^*: \mathbb{T}/Y \rightarrow \mathbb{T}/X$ has a left and a right adjoint $\sum_f \dashv f^* \dashv \prod_f$ called *pushforward along f* and *product along f* respectively. The product along f at u is the *object of local sections* of u , that is

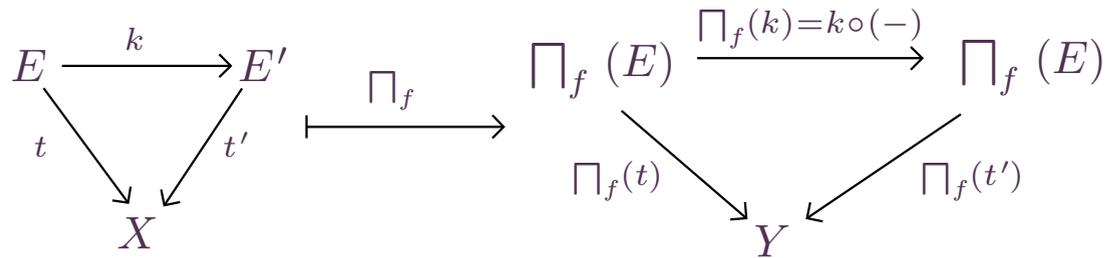
$$\prod_f (X \xrightarrow{u} Y) = \{ \sigma \in E^{f^{-1}(y)} \mid y \in Y, p \circ \sigma = i \}$$

in \mathbb{T} 's internal language.

Remark. A local section is thus given by the diagram



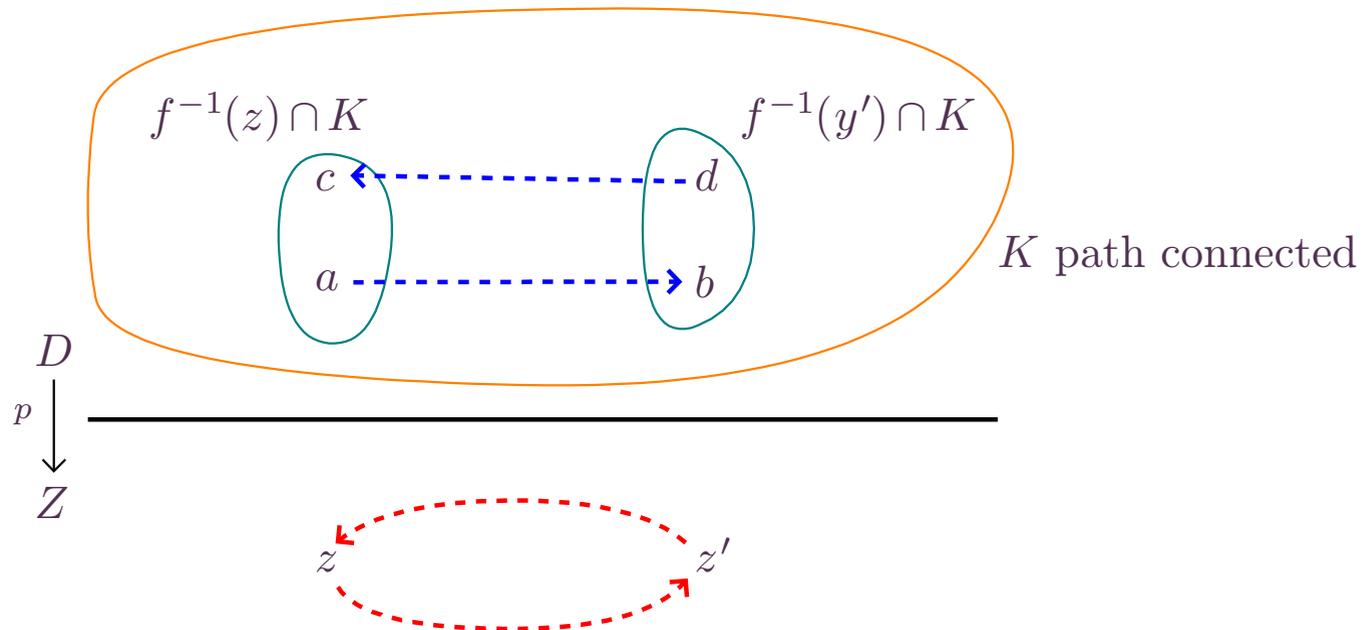
The action on maps is given by postcomposition



Lemma. Assume \mathbb{T} Hurewicz. Let $p: D \rightarrow Z$ be a fibration, $K \subseteq D$ a path connected component of D and $\delta: z \rightsquigarrow z'$ a path in Z . The following are equivalent

i. $f^{-1}(z) \cap K \neq \emptyset$;

ii. $f^{-1}(z') \cap K \neq \emptyset$

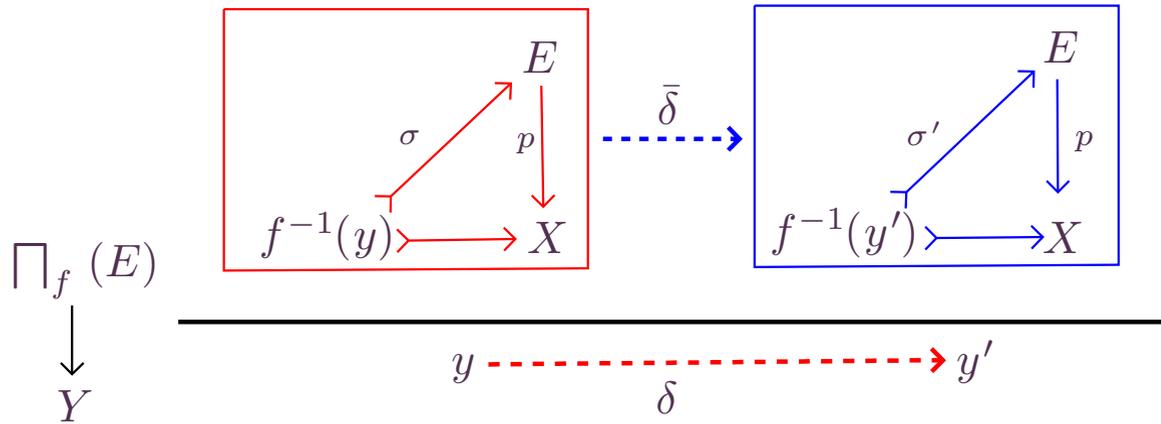


Theorem. *Let $p, f \in \mathbb{T}_1$ be fibrations. Then $\prod_f(p)$ is again a fibration.*

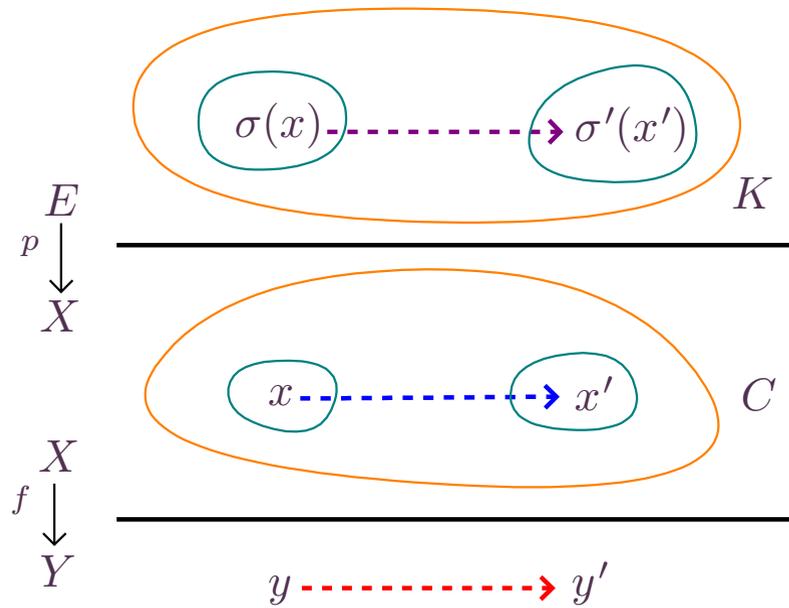
Definition. *Assume $X \in \mathbb{T}$. X 's local category is the subcategory $L_{\mathcal{F}}(X) \subseteq \mathbb{T}/X$ of the slice \mathbb{T}/X with all the objects being fibrations.*

Corollary. \prod_f restricts to local categories.

A path $\delta: y \rightsquigarrow y'$ in Y can be lifted to a path of sections:



since we have



The Hurewicz typos

Let \mathbb{C} be a category with finite limits. Recall that the product over A is given by pullback

$$(E, u) \times (E', u') = (E \times_A E', u * u')$$

with $u * u' = u \circ p_1 = u' \circ p_2$. The diagonal $\Delta_A E: (E, u) \rightarrow (E, u) \times (E, u)$ over A is thus given by

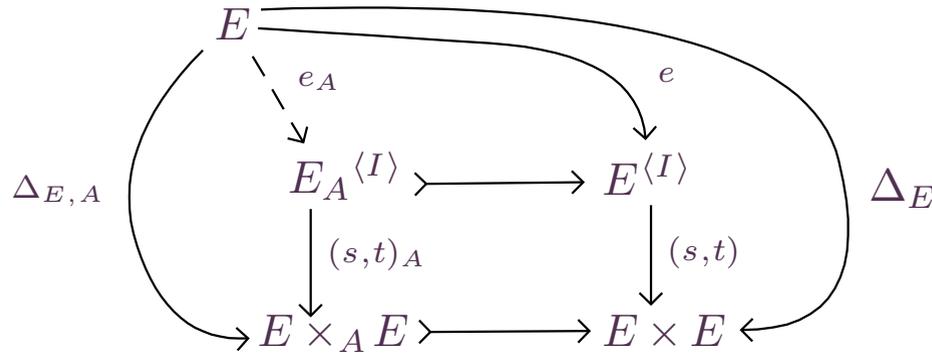
$$\Delta_{E,A} = (u * u) \circ (\text{id}_E, \text{id}_E)_A$$

Definition. Let $(\mathbb{J}, \mathcal{F})$ be a tribe. A fibration in \mathbb{J}/A is in \mathcal{F} (as a map over A).

Proposition. (Joyal) Let $(\mathbb{J}, \mathcal{F})$ be a tribe. The factorisation of the diagonal

$$\Delta_E = (s, t) \circ e$$

with e anodyne and (s, t) a fibration induces a factorisation in $L_{\mathcal{F}}(A)$ by pullback in \mathbb{J}



The path object is $(E_A^{<I>}, (u * u) \circ (s, t)_A)$. We have $s_A = p_1 \circ (u * u)$ and $t_A = p_2 \circ (u * u)$.

Lemma. Let \mathbb{T} be a Hurewicz topos and $f: A \rightarrow B$. The functorial square

$$\begin{array}{ccc} L_{\mathcal{H}}(A) & \xrightarrow{\square_f} & L_{\mathcal{H}}(B) \\ (-)_A^{\langle I \rangle} \downarrow & & \downarrow (-)_B^{\langle I \rangle} \\ L_{\mathcal{H}}(A) & \xrightarrow{\square_f} & L_{\mathcal{H}}(B) \end{array}$$

commutes.

Remark. This is much stronger than just preservation of the homotopy relation.

Theorem. *Any Hurewicz topos is a typos.*

